STA531 Midterm Exam 2 Solutions

- 1. Graphical models
 - (a) $p(x_1, \dots, x_8) = p(x_1)p(x_2)p(x_3)p(x_4|x_1, x_2)p(x_5|x_2, x_3)p(x_6|x_4)p(x_7|x_4, x_5)p(x_8|x_5)$
 - (b) Yes, Indeterminate, Yes, Indeterminate
 - (c)



- 2. Markov chains and graphical models
 - (a) Suppose p respect $G_{\rm D}$. Then

$$p(x_1, \dots, x_n) = p(x_1) \prod_{i=1}^{n-1} p(x_{i+1}|x_i) = \prod_{i=1}^{n-1} \varphi_i(x_i, x_{i+1})$$

if we define $\varphi_1(x_1, x_2) = p(x_1)p(x_2|x_1)$ and $\varphi_i(x_i, x_{i+1}) = p(x_{i+1}|x_i)$ for $i = 2, \ldots, n-1$. The functions $\varphi_1, \varphi_2, \ldots, \varphi_{n-1}$ are nonnegative, and the maximal cliques of G_U are $\{1, 2\}, \{2, 3\}, \ldots, \{n-1, n\}$. Therefore, p respects G_U .

(b) Suppose p respects $G_{\rm U}$. Then

$$p(x_1, \dots, x_n) = p(x_1) \prod_{i=2}^n p(x_i | x_1, \dots, x_{i-1}) = p(x_1) \prod_{i=2}^n p(x_i | x_{i-1}) = \prod_{i=1}^n p(x_i | x_{\operatorname{pa}(i)})$$

since

- $X_i \perp X_1, \ldots, X_{i-2} \mid X_{i-1}$ by the separation criterion for undirected graphical models, because all paths from i to $\{1, \ldots, i-2\}$ pass through i-1, and
- $pa(1) = \emptyset$ and $pa(i) = \{i 1\}$ for i = 2, ..., n.

Therefore, p respects $G_{\rm D}$.

3. Modeling the data collection process

As usual, the appropriate posterior on θ to use is $p(\theta|y_{\text{obs}}, I)$, since this is the posterior obtained by conditioning on all of the observations. Since $I_j = 1$ for all $j = 1, \ldots, n$, obs = $\{1, \ldots, n\}$. Further, the distribution of $I|y, \theta$ does not depend on θ (because $p(I|y, \theta) = \prod_{j=1}^{n} \phi_{y_j}^{I_j} (1 - \phi_{y_j})^{1-I_j} \propto_{\theta} 1$). Consequently, the data collection process plays no role:

$$p(\theta|y_{\text{obs}}, I) = p(\theta|y, I) \propto p(y, I, \theta) = p(\theta)p(y|\theta)p(I|y, \theta)$$
$$\propto_{\theta} p(\theta)p(y|\theta) \propto_{\theta} \theta^{a-1}(1-\theta)^{b-1} \prod_{j=1}^{n} \theta^{y_j}(1-\theta)^{1-y_j}$$
$$\propto_{\theta} \text{Beta}(\theta \mid a + \sum y_j, b + n - \sum y_j).$$

4. Viterbi algorithm

Define $\nu_n(z_n) = 1$ for each $z_n = 1, ..., m$, and denote $p(z_1) = p(z_1|z_0)$ for notational simplicity. Then

$$\max_{z_{1:n}} p(x_{1:n}, z_{1:n}) = \max_{z_{1:n}} \prod_{t=1}^{n} p(z_t | z_{t-1}) p(x_t | z_t)$$

$$\vdots$$

$$= \max_{z_{1:j}} \left(\prod_{t=1}^{j} p(z_t | z_{t-1}) p(x_t | z_t) \right) \underbrace{\max_{z_{j+1}} p(z_{j+1} | z_j) p(x_{j+1} | z_{j+1}) \nu_{j+1}(z_{j+1})}_{\text{call this } \nu_j(z_j)}$$

$$\vdots$$

$$= \max_{z_1} p(z_1) p(x_1 | z_1) \nu_1(z_1).$$

Therefore, we have the following algorithm:

- (1) Set $\nu_n(z_n) = 1$ for $z_n = 1, ..., m$.
- (2) For each j = n 1, n 2, ..., 1, for each $z_j = 1, ..., m$, compute

$$\nu_j(z_j) = \max_{z_{j+1}} p(z_{j+1}|z_j) p(x_{j+1}|z_{j+1}) \nu_{j+1}(z_{j+1}).$$

(3) Return $\max_{z_1} p(z_1) p(x_1|z_1) \nu_1(z_1)$.

(Note: If you showed that your algorithm takes nm^2 time, then that is great, but this was not required—as long as your algorithm actually does take nm^2 time, then that is sufficient.)